Generalized Loopy 2U: A New Algorithm for Approximate Inference in Credal Networks

<u>Alessandro Antonucci</u>, Marco Zaffalon, Sun Yi, Cassio de Campos IDSIA Lugano, Switzerland

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Models of uncertainty about the state of a categorical variable *X*

- A probability mass function P(X)
- More generally, a *closed convex set* of probability mass functions K(X) This is a *credal set* (Levi, 1980)
- Complete ignorance? A *vacuous* credal set $K_0(X)$
- Lower (and upper) expectation $E = \left[f(X) \right] = i \pi f \qquad \sum B(\pi) f(X)$

 $\underline{E}_{K}[f(X)] = \inf_{P(X) \in K(X)} \sum_{X} P(x)f(X)$

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Two main results

Theorem (about representation)

"Every credal net can be **equivalently represented** as a credal net over **binary** variables" (and the transformation takes only polynomial time)

Corollary (about inference)

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 - State of a variable as a joint state of a number of "bits"

 $X = x \iff (\tilde{X}^1 = \tilde{x}^1) \land (\tilde{X}^2 = \tilde{x}^2) \land \dots$

- Arcs binarization
 - For each arc between two variables, all the relative bits are linked
 - The bits of the same variable are completely connected

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 - Same calculations iterated over the conditional credal set $K(X_i|\pi_i)$ $\underline{P}(\tilde{x}_i^j|\mathrm{pa}(\tilde{X}_i^j)) = \mathrm{inf}_{P(X_i|\mathrm{pa}(X_i)) \in K(X_i|\mathrm{pa}(X_i))} P(\tilde{x}_i^j|\mathrm{pa}(\tilde{X}_i^j))$
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- Let $P(\mathbf{X})$ be the joint probability mass function of a BN,
- and $\tilde{P}(\tilde{X})$ the corresponding p.m.f. on the binarized BN.
- Then, $P(\mathbf{X}) = \tilde{P}(\tilde{\mathbf{X}})$.

CNs binarization is an outer approximation

- Let $K(\mathbf{X})$ be the strong extension of a CN,
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 the corresponding p.m.f. on the binarized BN.
- Then, $P(\mathbf{X}) = \tilde{P}(\tilde{\mathbf{X}})$.

CNs binarization is an outer approximation

- Let $K(\mathbf{X})$ be the strong extension of a CN,
- and $\tilde{K}(\tilde{\mathbf{X}})$ the strong extension of its binarization.
- Then, $K(\mathbf{X}) \subseteq \tilde{K}(\tilde{\mathbf{X}})$.

We can do better!

But another transformation should be applied before the binarization.

BNs binarization is exact

- Let $P(\mathbf{X})$ be the joint probability mass function of a BN,
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- For each i = 1,...,n, add a node T_i, parent of X_i, between X_i and pa(X_i)
- Decision nodes {T_i}ⁿ_{i=1} indexing the possible specifications of each mass function given the values of the parents
- Decision nodes can be regarded as chance nodes with a "vacuous" specification of the relative credal set
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- After DT specification, the nodes are either precise or vacuous
 - For precise specifications binarization is exact (result for BNs)
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- Binarization of DT-specified CNs is exact!
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- Multi-connected (non-binary) CNs?
- Binarization + L2U (twofold approx)
- DT + Bin + L2U = GL2U is better!
 - Approx only because of loopy!
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LUC Search	GL20	DIT+L20

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Multi-10			3.08
Multi-10			2.22
Multi-10			6.93
Multi-25			1.84
Multi-25			3.03
Polyt-50			2.89
Polyt-50			3.92
Insurance			1.75
Insurance			1.93
Alarm			3.02
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Multi-10		1.07	3.38
Multi-10		1.75	3.08
Multi-10		1.25	2.22
Multi-10		1.89	6.93
Multi-25		1.60	1.84
Multi-25		2.04	3.03
Polyt-50		1.93	2.89
Polyt-50		2.21	3.92
Insurance		1.17	1.75
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Multi-10	0.27	1.25	2.22
Multi-10	2.34	1.89	6.93
Multi-25	2.31	1.60	1.84
Multi-25	2.48	2.04	3.03
Polyt-50	1.12	1.93	2.89
Polyt-50	1.45	2.21	3.92
Insurance	0.55	1.17	1.75
Insurance	1.13	1.32	1.93
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- Exact binarization of BNs and CNs
- A state-of-the-art algorithm for CNs updating
- The algorithm of choice for very large nets?
- A Python/C++ implementation available (ask Sun Yi)
- Challenges
 - In numerical tests (G)L2U always converges.
 A formal proof of that?
 - For non-binary targets, accuracy can be improved with an alternative binarization of the target.

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